# Game Theoretic Approach for Supply Chain Optimization under Demand Uncertainty

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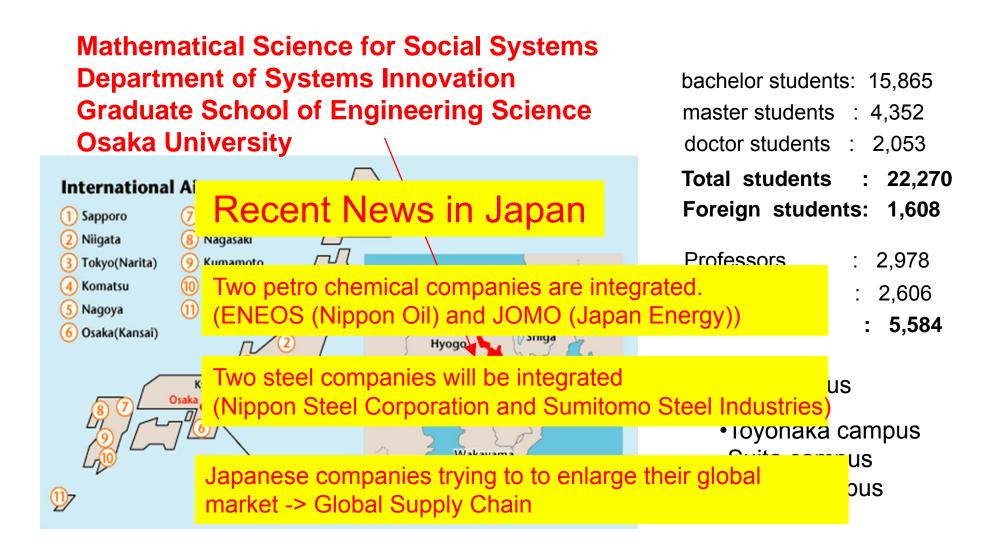
Mathematical Science for Social Systems Graduate School of Engineering Science Osaka University March 9, 2011

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# Agenda

- 1. Introduction of our group
- 2. Why game theoretic approach is required ?
- 3. Nash Equilibrium
- 4. Stackelberg game
- 5. Supply chain coordination with quantity discount contract
- 6. Numerical example
- 7. Conclusion and future works

# **Osaka University**



# Collaboration with Japanese industries

Semiconductor factory automation

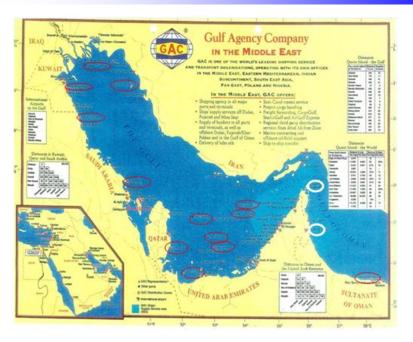
Control of automated guided vehicles for transportation Scheduling of cluster tool for silicon wafer production

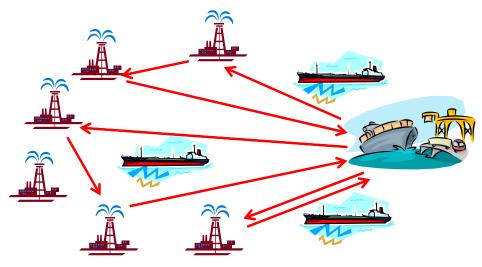
• Railway scheduling automation

railway crew scheduling in Japan Railway Train-set scheduling with maintenance constraints Shift scheduling

Petroleum chemical industry automation
 International ship scheduling for crude oil transportation
 Transportation network design

# International ship scheduling problem



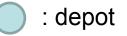


loading places

Split Delivery Vehicle Routing Problem (SDVRP) Column generation is applied to solve the problem efficiently. Is it OK for practitioners ? No, there will be uncertainty Human operators are negotiating timing, due dates and costs in practical.

# Transportation network design





• : city

Energy demand will be change oil to electricity

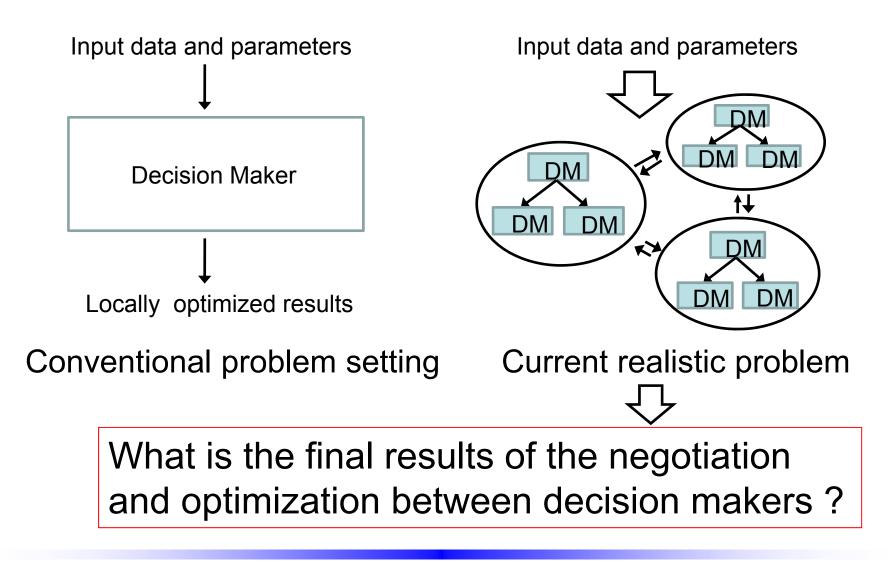
How many number of depots will be required to satisfy demands from customers to minimize distribution costs ?

Classical facility location problem

Is it OK for practitioners? The answer is no. There will be competing companies. We have to consider future marketing and competition.

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## Why game theoretic approach is required ?



## Game theoretic approach

simultaneous optimization > game theoretic approach > separate optimization

#### Game theoretical analysis

- Stackelberg game on channel coordination (Cai et al. 2009)
  - leader and follower

#### Nash equilibrium analysis

- Equilibrium behavior of decentralized supply chain with competing retailers (Bernstein et al. 2005) —analysis on supermodularity of profit function
- Price competing model with risky supplier (Serel, 2008)
   —analysis on quasi-concavity of profit function

## Nash equilibrium

For n player game,

 $\mathcal{X}_i$  is the strategy for player  $\, m{i}$  ,  $\mathcal{X}_{-i}$  is the strategy for players except player  $\, m{i}$  .

The strategy  $(x_1^*, x_2^*, \dots, x_n^*)$  is Nash equilibrium when the optimal strategy for player i is  $x_i^*$ on the condition that all the players except player i is  $x_{-i}^*$ 

$$x_i^* = \arg \max_{x_i} \pi_i(x_i, x_{-i}^*)$$

## Supermodularity

The real-valued function f(x) is supermodular if

 $f(x') + f(x'') \le f(\min\{x', x''\}) + f(\max\{x', x''\})$ 

The real-valued function g(x,t) is increasing differences if  $g(x',t'') - g(x',t') \le g(x'',t'') - g(x'',t')$ 

The 2<sup>nd</sup> differential function f(x) is increasing differences if

$$\frac{\partial f(x)}{\partial x_i \partial x_j} \ge 0$$

## Supermodularity and Nash equilibrium (Topkis 1998)

#### Theorem 1

f(x) is supermodular when f(x) is increasing differences and f(x) is supermodular on the fixed  $X_i$ 

#### Theorem 2

If the profit function  $\pi_i$  is supermodular,

Nash equilibrium exists when player i determines the strategy by

$$x_i^{k+1} = \arg\max_{x_i} \pi_i(x_i, x_{-i})$$

Existence of Nash equilibrium

Supermodular function  $\phi(\max\{x, y\}) + \phi(\min\{x, y\}) \ge \phi(x) + \phi(y)$ x, y: pricing vector

The profit function for each company is supermodular  $\downarrow$ Increase or decrease of pricing is identical  $\downarrow$ Existence of Nash equilibrium Game theoretic approach for supply chain coordination under demand uncertainty with quantity discount contract

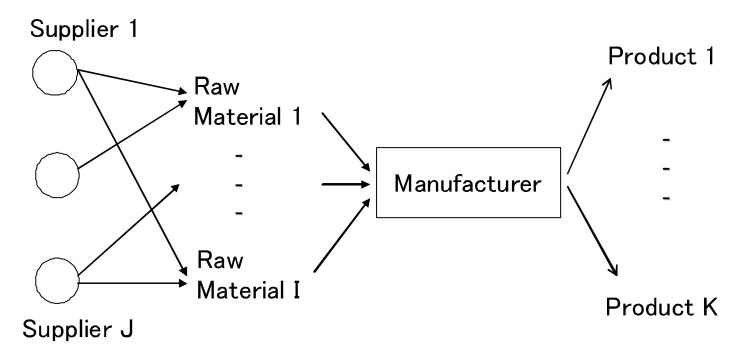
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# Outline

- 1. Game theoretic approach
- 2. Problem definition
- 3. suppliers and manufacturers model
- 4. Stackelberg game
- 5. Quantity discount as a coordination mechanism
- 6. Optimization algorithm
- 7. Numerical Example

# **Problem definition**



A supply network consists of a manufacturer and its suppliers

- Multiple products
- Uncertain demand
- Multiple suppliers
- Capacity limits

# Problem

How to select suppliers for each raw material/part and determine purchasing allocation of each RM among the suppliers to maximize the total profit of the manufacturer and suppliers

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Manufacturer (leader)	<ul> <li>Simultaneously determine</li> <li>optimal annual production levels</li> <li>supplier selection</li> <li>demand uncertainty</li> </ul>
suppliers (followers)	<ul> <li>quantity of raw material</li> <li>price of raw material</li> <li>setup and ordering cost</li> </ul>

## Literature review

#### **Quantity discount model:**

Kim et al.(2002) Configuring a Manufacturing Firm's Supply Network with Multiple Suppliers

Tsai (2007) An optimization Approach for Supply Chain Management Models with Quantity Discount Policy

Zhang and Ma (2009) Optimal Acquisition Policy with Quantity Discounts and Uncertain Demands

#### **Coordination Model:**

Qin et al. (2007) Channel Coordination and Volume Discounts with Price-sensitive Demand

Yu et al. (2009) Stackelberg Game-theoretic Model for Optimizing Advertising ,Pricing and Inventory Policies in Vendor Managed Inventory Production Supply Chains

Leng et al. (2010) Game-theoretic Analyses of Decentralized Assembly Supply Chains

## Modeling: assumptions

•A manufacturer and suppliers two-tier supply chain (leader-follower relationship)

 One cycle of the manufacturer's long term production period

- •PDF f(z) of demand z is known
- •Unit understocking cost  $a_k$ , and unit overstocking cost  $b_k$  are known

## Decision variables and parameters

#### Decision variables

 $\begin{array}{ll} D_{ij} & \mbox{quantity of raw materials i purchased from supplier j.} \\ y_k & \mbox{production quantity of product k for manufacturer.} \\ w_j \in \{0,1\} & \mbox{binary variable indicating whether manufacturer buys from supplier j.} \end{array}$ 

#### Parameters

 $z_k$  random demand for product k on normal distribution

Maximize 
$$J = \sum_{k=1}^{K} \{\int_{0}^{y_{k}} [r_{k}z_{k} - b_{k}(y_{k} - z_{k})f(z_{k})dz_{k} + \int_{y_{k}}^{\infty} r_{k}y_{k} - a_{k}(z_{k} - y_{k})f(z_{k})dz_{k}\} - \sum_{k=1}^{K} e_{k}y_{k} - \sum_{i=1}^{I} \sum_{j=1}^{J} q_{ij} D_{ij} - \sum_{j=1}^{J} m_{j} w_{j}$$

<u>Total Profit</u>=<u>Expected Revenue</u> (while it is overstocking and understocking)-<u>Production</u> <u>cost-Procurement cost-Management cost</u>

Subject to

$$\begin{split} & \sum_{k=1}^{K} y_k g_k \leq \sum_{j=1}^{J} D_{ij}, \forall i = 1, \dots I \quad (1) \\ & \text{Bills of Materials} \\ & \sum_{k=1}^{K} t_k y_k \leq C \quad (2) \\ & \text{Production Capacity} \\ & \sum_{i=1}^{I} n_{ij} D_{ij} \leq c_j w_j, \forall j = 1, \dots J \quad (3) \\ & \text{Procurement} \\ & \text{Capacity} \\ & D_{ij}, y_k \geq 0, w_j \in \{0,1\}, \forall i, \forall j, \forall k \quad (4) \\ & \text{Variables} \end{split}$$

## Supplier's EOQ model

#### Total profit = Management cost

+Gross revenue- Inventory holding cost –Order processing cost

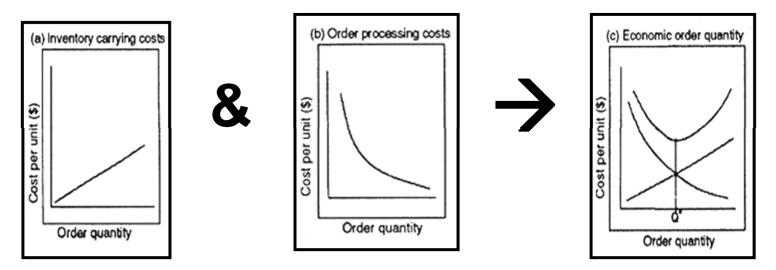
(5)

(6)

Maximize 
$$\Pi_j = m_j w_j + \sum_{i=1}^{I} [(q_{ij} - h_{ij}) D_{ij} - \frac{A_{ij} Q_{ij}}{2} - \frac{S_{ij} D_{ij}}{Q_{ij}}]$$

Subject to

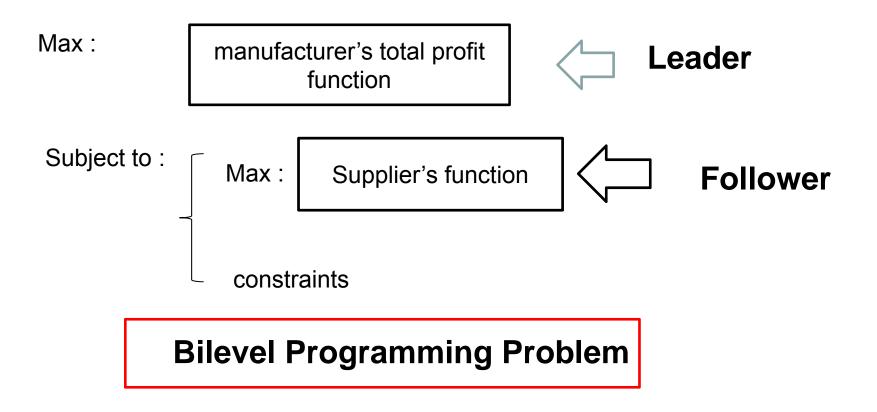
$$q_{ij} \geq h_{ij}, Q_{ij} > 0, \forall i = 1, \dots, I, \forall j = 1, \dots, J$$



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## Stackerberg game model

#### Stackerberg Game model

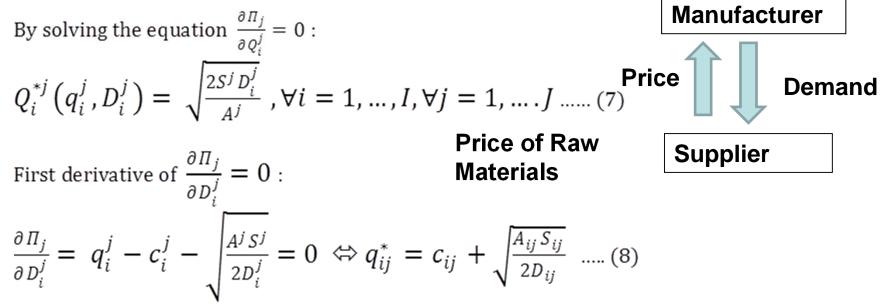


## Supplier's optimal response function

#### The supplier's optimal response function

First derivative of supplier's profit function (5) with respect to  $Q_i$ :

$$\frac{\partial \Pi_j}{\partial Q_i^j} = -\frac{S^j D_i^j}{Q_i^{j^2}} - \frac{A^j}{2}, \forall i = 1, \dots, I, \forall j = 1, \dots, J$$

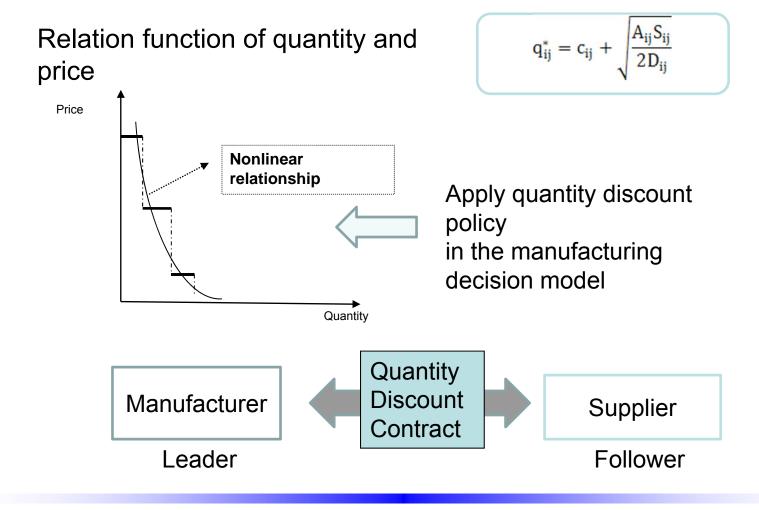


**Quantity of purchased materials** 

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## Coordination between manufacturer and suppliers

#### Analysis of the Stackelberg game equilibrium



# Manufacturing planning model with quantity discount model

The quantity discount model can be applied to solve manufacturer's decision problem

<u>Constrains:</u> (1), (2), (3), (4), (5), (6)

$$\begin{array}{ll} D_{ijl} \leq d_{ijl}^{H} u_{ijl}, \forall i, j, l \\ D_{ijl} \geq d_{ijl}^{S} u_{ijl}, \forall i, j, l \\ \sum_{l=1}^{L_{j}} u_{ijl} = v_{ij}, \forall j \\ w_{j} \geq v_{ij}, \forall \\ u_{ijl}, v_{ijl}, w_{j} \in \{0, 1\}, \forall i, j, l \end{array}$$

$$(10)$$

$$(11)$$

$$(12)$$

$$(13)$$

$$(14)$$

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## **Optimization algorithm**

Fix 0-1 binary variables  $u_{ijl}$ ,  $v_{ij}$ ,  $w_j$  to a 0-1 binary configuration

#### Primal Problem:

$$\max \sum_{k=1}^{K} \left\{ \int_{0}^{y_{k}} [r_{k}z_{k} - b_{k} (y_{k} - z_{k})] f(z_{k}) dz_{k} + \int_{y_{k}}^{\infty} [r_{k}y_{k} - a_{k}(z_{k} - y_{k})] f(z_{k}) dz_{k} \right\} \\ -\sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{l=1}^{L_{j}} c_{ijl} x_{ijl} - \sum_{k=1}^{K} e_{k} y_{k}$$
solved by SQP algorithm

#### Master Problem:

$$\min(\mu OA + \sum_{j=1}^{J} (\sum_{i=1}^{I} c_{ij} x_{ij}) \sum_{l=1}^{L_j} (1 - d_{jl}) u_{jl} + \sum_{j=1}^{J} m_j w_j)$$

s.t 
$$\mu OA \ge -P(y_k^h) - \frac{\delta P}{\delta y_k}\Big|_{y_k = y_k^h} (y_k - y_k^h)$$

$$\sum_{j \in B^h} (w_j + \sum_{i \in B^h} v_{ij} + \sum_{l \in B^h} u_{jl}) - \sum_{j \in NB^h} (w_j + \sum_{i \in NB^h} v_{ij} + \sum_{l \in NB^h} u_{jl}) \le |B^h| - 1$$

where 
$$P(y) = -\sum_{k=1}^{K} \left\{ \int_{0}^{y_{k}} [r_{k}z_{k} - b_{k}(y_{k} - z_{k})]f(z_{k})dz_{k} + \int_{y_{k}}^{\infty} [r_{k}y_{k} - a_{k}(z_{k} - y_{k})]f(z_{k})dz_{k} \right\} - \sum_{k=1}^{K} e_{k}y_{k}$$

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#### Improvement of OA algorithm

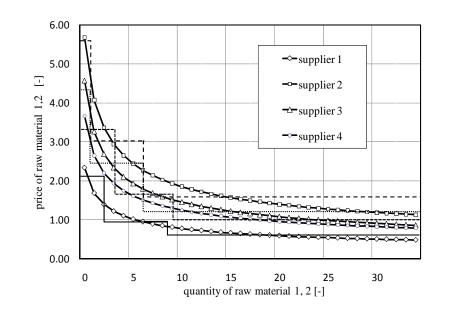
The deviation between realized demand and production.

$$\begin{array}{l} \min\left(z_{k}, y_{k}\right) & \text{Set} & \left(x_{k} = \frac{z_{k} - \tilde{z}}{\sigma_{k}}\right) \left(Y_{k} = \frac{y_{k} - \tilde{z}}{\sigma_{k}}\right) \\ \text{Therefore, the expected revenue from products is } \mathsf{E}(RE_{k}) = \Phi(Y_{k})\mathsf{E}[r_{k}z_{k}] \left[\frac{z_{k} - z_{k}}{\sigma_{k}} \le \frac{y_{k} - \tilde{z}}{\sigma_{k}}\right] \\ RE_{k}(Z_{k}) = r_{k} \min\left(z_{k}, y_{k}\right) = \begin{cases} r_{k}z_{k} & \text{if } z_{k} < y_{k} \\ r_{k}y_{k} & \text{if } z_{k} > y_{k} \end{cases} + (1 - \Phi(Y_{k})) \mathsf{E}[r_{k}y_{k}] \left[\frac{z_{k} - z}{\sigma_{k}} \le \frac{y_{k} - \tilde{z}}{\sigma_{k}}\right] \\ \text{According to the definition of normal distribution:} \\ = r_{k} \tilde{z} + r_{k}\sigma_{k}\{\Phi(Y_{k})\mathsf{E}[X_{k}|X_{k} \le Y_{k}]\} \\ \mathsf{F}(x) = \frac{1}{\sqrt{2\pi}} e^{\left(-\frac{1}{2}(x)^{2}\right)} & \mathsf{E}[Y_{k}|Y_{k} \le X_{k}] = Y_{k} \\ = \Rightarrow \Phi(z) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{z} e^{-\frac{1}{2}t^{2}} dt \\ \mathsf{Final Equation:} & \mathsf{E}[X_{k}|X_{k} \le Y_{k}] = \frac{\sqrt{\frac{1}{2\pi}}\int_{-\infty}^{Y_{k}} x_{k}e^{-\frac{1}{2}x_{k}^{2}} dx_{k}}{\frac{1}{\sqrt{2\pi}}\int_{-\infty}^{z} e^{-\frac{1}{2}x_{k}^{2}} dx_{k}} = \frac{-1}{\Phi(Y_{k})} \frac{e^{-\frac{1}{2}x_{k}^{2}}}{\sqrt{2\pi}} = -\frac{F(Y_{k})}{\Phi(Y_{k})} \\ \mathsf{E}(\text{overproduction}) = b_{k}\{\tilde{z} + \sigma_{k}[-\mathsf{F}(Y_{k}) + (1 - \Phi(Y_{k})) Y_{k}]\} \\ \mathsf{E}(\text{shortfall}) = a_{k}\{\mathsf{E}\{\min(z_{k}, y_{k})\} - \check{z}\} \\ = a_{k} \sigma_{k}[-\mathsf{F}(Y_{k}) + (1 - \Phi(Y_{k})) Y_{k}] \\ \mathsf{E}(RE_{k}) = \check{z} + \sigma_{k}[-\mathsf{F}(Y_{k}) + (1 - \Phi(Y_{k})) Y_{k}] \end{aligned}$$

## Numerical example

4 suppliers, 1 manufacturer, 2 finished products,5 raw material,3 quantity discount interval

An Intel(R) Core2Duo E8400 3 GHz with 3GB memory (Matlab R2009a and CPLEX 9.0 with CPLEXINT library)





## Numerical example

Optimal Solu	suppl.1	
		mater.1
Revenue	28,169	mater.2
Overstock cost	5.2	mater.3
Shortage	1,378	mater.4
cost	,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	mater.5
Production cost	5,090	suppl.2
Material cost	98.2	mater.1-5
Management	259	suppl.3
cost		mater.1-5
Total profit	21,338.6	suppl.4
Manufacturer's	mater.1	

Results: Total computation time: 28.9 second The derived upper bound is 21365.6, and the lower bound is 21338.6	
Duality gap : 0.127%	

$\frown$						
suppl.1	$D_{i11}$	$D_{i12}$	$D_{i13}$	$q_{i1}^{*}$	$Q_{i1}^*$	EOQi1
mater.1	0	0	29.48	0.51	12.1	-0.412
mater.2	0	0	27.38	0.53	11.7	-0.427
mater.3	0	0	13.00	0.59	5.10	-0.392
mater.4	0	0	21.50	0.50	6.56	-0.305
mater.5	0	0	33.64	0.44	8.20	-0.244
suppl.2	$D_{i21}$	D <sub>i22</sub>	D <sub>i23</sub>	$q_{i2}^{*}$	$Q_{i2}^*$	EOQi2
mater.1-5	0	0	0	-	-	0
suppl.3	$D_{i31}$	D <sub>i32</sub>	$D_{i23}$	<b>q</b> <sup>*</sup> <sub>i3</sub>	$Q_{i3}^*$	EOQi3
mater.1-5	0	0	0	-	-	0
suppl.4	$D_{i41}$	$D_{i42}$	$D_{i43}$	$q^*_{i4}$	$Q_{i4}^*$	EOQi4
mater.1						
mater. I	0	0	0	-	-	0
mater.2	0 0	0 3.50	0 0	- 2.05	- 4.32	0 -1.852
	-	_	-			-
mater.2	0	3.50	0	2.05	4.32	-1.852

Supplier's Optimal Solution

Contract decision models including quantity discount model and the Stackelberg game theoretic model are studied.

The outer approximation method is applied to solve the mixed integer nonlinear programming problem.

The computational experiments demonstrate the effectiveness of the proposed method and the feasibility of the model.

## Conclusion and future works

#### In the Future

#### Contract Decision for Supply Chain Optimization

Integrated supply chain management is investigated to maximize the profits, meanwhile, to decrease the energy consumption and the burden of environment

#### International Logistics

Transportation, distribution, manufacturing and the like to support the globalization of the business

#### Transportation Network Design

Reasonable routing with consideration of the future market and trade off